A Journey in Functional Programming An introduction to Haskell

Davide Spataro¹

¹Department of mathematics And Computer Science University of Calabria

October 14, 2015

Table of contents I

Introduction - Functional Programming

Scope Of The Talk Functional Programming Tools and Installation Toolbox - Hello world(s)

Basics - Syntax Arithmetic And Boolean algebra Guards, where, let if and case construct Ranges List Lambda Functions

Coding - Problems on Lists

Last element kth element Palindrome List

Table of contents II

Problem on Numbers

Primality Test Greatest common divisor Euler's torient

Find Best Variance - Stock Data I/O - Find Best Variance

Coding - Project Euler Problem 1 Problems 1

Coding - Project Euler Problem 26 Problems 26

Section 1

Introduction - Functional Programming

How many of you are capable of writing (correct) quicksort in an imperative language at the first try?

All Haskell programmers are!

What does this code do?

```
void function (int *a, int n) {
    int i, j, p, t;
    if (n < 2)
        return;
    p = a[n / 2];
    for (i = 0, j = n - 1;; i++, j--) {
        while (a[i] < p)</pre>
           i++;
        while (p < a[j])
           i--;
        if (i \ge j)
           break;
        t = a[i];
        a[i] = a[j];
        a[j] = t;
    }
    function(a, i);
    function(a + i, n - i);
}
```

```
function ::(Ord a) => [a] -> [a]
function [] = []
function (x:xs) = (function 1) ++ [x] ++ (function g)
    where
    l = filter (<x) xs
    g = filter (>=x) xs
```

- No indices
- No memory/pointer management
- No variable assignment

Functional Programming

Definition and Intuitive idea

- Computation is just function evaluation
 ≠ program state manipulation.
- ► Based on λ-calculus that is an alternative (to set theory) and convenient formalization of logic and mathematics for expressing computation
- Logic deduction ⇔ λ−calculus thanks to the Curry-Howard correnspondence.
- A program is a proof!
- Heavily based on Category Theory (Monad, Functor, etc.)



Figure: Alonzo Church, father of λ -calculus

Imperative vs Functional

Imperative

- Focus on low-level how!
- A program is an ordered sequence of instructions
- Modifies/track the program's state
- Functional
 - Focus on High level what!
 - Specify high-level transformation/constraint on the desidered result description.

Functional

sum [1..99]

```
Imperative, suffer from the so called
indexitis
unsigned int sum=0;
for(int i=1;i<100;i++)
sum+=i;
```

Characteristic	Imperative	Functional
Programmer focus	Algorithm design	What the output look like?
State changes	Fundamental	Non-existent
Order of execution	Important	Low importance
		(compilers may do much work on this)
Primary flow control	Loops, conditionals	Recursion and Functions
Primary data unit	Structures or classes	Functions

► Other pure/quasi-pure languages: Erlang, Scala, F, LISP.

- 1. Haskell's expressive power can improve productivity/understandability/maintanibility
 - Get best from compiled and interpreted languages
 - Can understand what complex library does
- 2. Strong typed Catches bugs at compile time
- 3. Powerful type inference engine i.e. **no need to explicitely specify types**
- New Testing metologies. Proving vs Testing (e.g. QuickCheck)
- 5. Automatic parallelization due to code purity.

What really is Haskell?

Purely Functional language

- Functions are first-class object (same things as data)
- ► Deterministic No Side Effect- same function call ⇒ same Ouput, EVER!
- safely replace expressions by its (unique) result value This referial transparency leaves room for compiler optimization and allow to mathematically prove correctness
- Evaluate expression rather than execute instruction
- Function describes what data are, not what what to do to...
- Everything is immutable (i.e. NO variables)
- Multi-parameters function simply does not exists.

It won't execute anything until is *really* needed

- It is possible to define and work with infinite data structures
- Define new control structure just by defining a function.
- Reasoning about time/space complexity much more complicated



Understanding laziness

```
lazyEval 0 b = 1 
lazyEval _ b = b
```

b never computed if the first parameter is zero

Strict evaluation: parameter are evaluated **before** to be passed to functions

```
1
2
3
4
5
6
```

```
int cont=0;
auto fcall = [] (int a, int b)
{if(a==0) return 1; else return b;};
auto f1 = [] () { cont++; return 1};
auto f2 = [] () { cont+=10; return 2};
fcall (f1(),f2()));
```

How many times cont is updated?

Understanding laziness

```
lazyEval 0 b = 1 
lazyEval _ b = b
```

b never computed if the first parameter is zero

Strict evaluation: parameter are evaluated **before** to be passed to functions

```
1
2
3
4
5
6
```

```
int cont=0;
auto fcall = [] (int a, int b)
{if(a==0) return 1; else return b;};
auto f1 = [] () { cont++; return 1};
auto f2 = [] () { cont+=10; return 2};
fcall (f1(),f2()));
```

How many times cont is updated? ALWAYS twice

Problem: compute the k^{th} Fibonacci number.

Problem: compute the k^{th} Fibonacci number. f a b k = if k==0 then a else f b (a+b) (k-1)

▶ Defines a recursive function *f* that takes *a*,*b*,*k* as parameters.

Problem: compute the k^{th} Fibonacci number.

- ▶ Defines a recursive function *f* that takes *a*,*b*,*k* as parameters.
- Spaces are important. Are like function call operator () in C-like languages.

Problem: compute the k^{th} Fibonacci number.

- ▶ Defines a recursive function *f* that takes *a*,*b*,*k* as parameters.
- Spaces are important. Are like function call operator () in C-like languages.
- ► Wait, three space in f a b k: 3 function calls? YES!. Every function in Haskell officially only takes one parameter.

Problem: compute the k^{th} Fibonacci number.

- ▶ Defines a recursive function *f* that takes *a*,*b*,*k* as parameters.
- Spaces are important. Are like function call operator () in C-like languages.
- ► Wait, three space in f a b k: 3 function calls? YES!. Every function in Haskell officially only takes one parameter.
- f infact has type
- f :: Integer ->(Integer ->(Integer ->Integer))
 i.e. a function that takes an integer and return (the ->) a
 function that takes an integer and returns ...

Problem: compute the k^{th} Fibonacci number.

- ▶ Defines a recursive function *f* that takes *a*,*b*,*k* as parameters.
- Spaces are important. Are like function call operator () in C-like languages.
- ► Wait, three space in f a b k: 3 function calls? YES!. Every function in Haskell officially only takes one parameter.
- f infact has type
- f :: Integer ->(Integer ->(Integer ->Integer))
 i.e. a function that takes an integer and return (the ->) a
 function that takes an integer and returns ...
- f 0 :: Integer ->(Integer ->Integer)

Problem: compute the k^{th} Fibonacci number.

f a b k = if k==0 then a else f b (a+b) (k-1)

- ▶ Defines a recursive function *f* that takes *a*,*b*,*k* as parameters.
- Spaces are important. Are like function call operator () in C-like languages.
- ► Wait, three space in f a b k: 3 function calls? YES!. Every function in Haskell officially only takes one parameter.
- f infact has type
- f :: Integer ->(Integer ->(Integer ->Integer))
 i.e. a function that takes an integer and return (the ->) a
 function that takes an integer and returns ...

```
f 0 :: Integer ->(Integer ->Integer)
```

f 0 1 :: Integer->Integer

Problem: compute the k^{th} Fibonacci number.

- ▶ Defines a recursive function *f* that takes *a*,*b*,*k* as parameters.
- Spaces are important. Are like function call operator () in C-like languages.
- ► Wait, three space in f a b k: 3 function calls? YES!. Every function in Haskell officially only takes one parameter.
- f infact has type
- f :: Integer ->(Integer ->Integer))
 i.e. a function that takes an integer and return (the ->) a
 function that takes an integer and returns ...
 f 0 :: Integer ->(Integer ->Integer)

```
f 0 1 :: Integer->Integer
```

```
f 0 1 10 :: Integer
```

Currying directly and naturally address the high-order functions support Haskell machinery.

High-order function:

- Takes functions as parameter
- "returns" a function

Currying directly and naturally address the high-order functions support Haskell machinery.

High-order function:

- Takes functions as parameter
- "returns" a function

zipwith

- Combines two list of type a and b using a function f that takes a parameter of type a and one of type b and return a value of type c, producing a list of elements of type c.
- ▶ zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]

$zipWith :: (a \rightarrow b \rightarrow c) \rightarrow [a] \rightarrow [b] \rightarrow [c]$

zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith _ [] = []

zipWith :: (a -> b -> c) -> [a] -> [b] -> [c] zipWith _ [] = [] zipWith _ [] _ = []

zipWith :: (a -> b -> c) -> [a] -> [b] -> [c] zipWith _ [] = [] zipWith _ [] _ = [] zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys

zipWith :: (a -> b -> c) -> [a] -> [b] -> [c] zipWith _ [] = [] zipWith _ [] _ = [] zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys usage examples #> zipWith (+) [1,2,3] [4,5,6] = [5,7,9]

```
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith _ [] = []
zipWith _ [] _ = []
zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys
usage examples
#> zipWith (+) [1,2,3] [4,5,6] = [5,7,9]
#> zipWith (*) [1,2,3] [4,5,6] = [4,10,18]
```

```
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith _ [] = []
zipWith _ [] _ = []
zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys
usage examples
#> zipWith (+) [1,2,3] [4,5,6] = [5,7,9]
#> zipWith (*) [1,2,3] [4,5,6] = [4,10,18]
#> let f = in (\a b -> (ord a) + b)
zipWith f ['A'..] [1..]
```

```
zipWith :: (a \rightarrow b \rightarrow c) \rightarrow [a] \rightarrow [b] \rightarrow [c]
zipWith _ _ [] = []
zipWith _ [] _ = []
zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys
usage examples
\# > zipWith (+) [1,2,3] [4,5,6] = [5,7,9]
\# > \text{zipWith} (*) [1,2,3] [4,5,6] = [4,10,18]
\# let f = in (\a b -> (ord a) + b)
             zipWith f ['A'..] [1..]
What about this call? (missing one parameter)
let l = zipWith (*) [1,2,3]
1 is a function that takes ONLY a list of Integer and returns
[1a, 2b, 3 dotc]
```

- Haskell is stricly typed
- Helps in thinking and express program structure
- Turns run-time errors into compile-time errors. "If it compiles, it must be correct", is moslty true.

Abstraction

Every idea, algorithm, and piece of data should occur exactly once in your code. Haskell features as **parametric polymorphis**¹, **typeclasses**² **high-order functions** greatly aid in fighting repetition.

¹templates with type inference ²Interfaces

```
C-like vs Haskell
Code as the one that follows
int acc = 0;
for ( int i = 0; i < lst.length; i++ )
    acc = acc + 3 * lst[i];
is full of low-level details of iterating over an array by keeping track
of a current index. It much elegantely translates in:
    sum (map (*3) lst)</pre>
```

```
Other examples:

partition (even) [49, 58, 76, 82, 83, 90]

--prime number generation

let pgen (p:xs) = p : pgen [x|x <- xs, x'mod'p > 0]

take 40 (pgen [2..])
```
A full comprehensive, development environment for Haskell³⁴. Installation

\$sudo apt-get install haskell-platform

GHC (Great Glasgow Compiler): State of the art
 GHCi A read-eval-print loop interpreter
 Cabal Build/distribuite/retrieve libraries
 Haddock A high quality documentation generation tool for
 Haskell

³https://www.haskell.org/platform/index.html
⁴http://tryhaskell.org/

Our First Program

```
Create a file hello.hs and compile with the followings
main = putStrLn "Hello World with Haskell"
$ghc -o hello hello.hs
```

GHCi

```
Execute and play with GHCi by simply typing
reverse [1..10]
#> :t foldl
#> [1..]
#> (filter (even) .reverse) [1..100]
```

Section 2

Basics - Syntax

Syntax Basics

Arithmetic and Boolean algebra works as expected v1 = 12 v2 = mod (v1+3) 10 v3 = not \$ True || (v2>=v1) --not (True || (v2>=v1))

Syntax Basics

- Arithmetic and Boolean algebra works as expected v1 = 12 v2 = mod (v1+3) 10
 - v3 = not \$ True || (v2>=v1) --not (True || (v2>=v1))
- Function definition is made up of two part: type and body. The body is made up of several *clause* that are evaluated (pattern matched) top to bottom.

```
4 exp _ 0 = 1

5 exp 0 _ = 0

6 exp a b = a * (exp a (b-1))

What if we swap line 2 and 3?
```

Syntax Basics

- Arithmetic and Boolean algebra works as expected
 v1 = 12
 v2 = mod (v1+3) 10
 - v3 = not \$ True || (v2>=v1) --not (True || (v2>=v1))
- Function definition is made up of two part: type and body. The body is made up of several *clause* that are evaluated (pattern matched) **top to bottom**.

```
7 exp _ 0 = 1
8 exp 0 _ = 0
9 exp a b = a * (exp a (b-1))
What if we swap line 2 and 3?
```

Comments:

```
--this is an inline comment
{-
All in here is comment
-}
```

```
Guards,let and where constructs
  fastExp :: Integer -> Integer-> Integer
1
2 fastExp _0 = 1
3 fastExp a 1 = a
4
  fastExp a b
5
     |b < 0 = undefined
6
     leven b = res*res
7
     |otherwise = let next=(fastExp a (b-1)) in (a * next)
8
      where res=(fastExp a (div b 2))
  Suppose we execute fastExp 2 7. The call stack would be
    fastExp 2 7 line 7 pattern match
```

⁵Here for more informations: https://wiki.haskell.org/Let_vs_Where

```
Guards,let and where constructs
  fastExp :: Integer -> Integer-> Integer
1
2 fastExp _0 = 1
3 fastExp a 1 = a
4
  fastExp a b
5
     |b < 0 = undefined
6
    |even b = res*res
7
     |otherwise = let next=(fastExp a (b-1)) in (a * next)
8
      where res=(fastExp a (div b 2))
  Suppose we execute fastExp 2 7. The call stack would be
    fastExp 2 7 line 7 pattern match
    fastExp 2 6 line 6 pattern match
```

⁵Here for more informations: https://wiki.haskell.org/Let_vs_Where

```
Guards,let and where constructs
  fastExp :: Integer -> Integer-> Integer
1
2 fastExp _0 = 1
3 fastExp a 1 = a
4
  fastExp a b
5
     |b < 0 = undefined
6
     leven b = res*res
7
     |otherwise = let next=(fastExp a (b-1)) in (a * next)
8
      where res=(fastExp a (div b 2))
  Suppose we execute fastExp 2 7. The call stack would be
     fastExp 2 7 line 7 pattern match
     fastExp 2 6 line 6 pattern match
     fastExp 2 3 line 7 pattern match
```

⁵Here for more informations: https://wiki.haskell.org/Let_vs_Where

```
Guards,let and where constructs
  fastExp :: Integer -> Integer-> Integer
1
2 fastExp _0 = 1
3 fastExp a 1 = a
4
  fastExp a b
5
     |b < 0 = undefined
6
     leven b = res*res
7
     |otherwise = let next=(fastExp a (b-1)) in (a * next)
8
      where res=(fastExp a (div b 2))
  Suppose we execute fastExp 2 7. The call stack would be
     fastExp 2 7 line 7 pattern match
     fastExp 2 6 line 6 pattern match
     fastExp 2 3 line 7 pattern match
     fastExp 2 2 line 6 pattern match
```

⁵Here for more informations: https://wiki.haskell.org/Let_vs_Where

```
Guards,let and where constructs
  fastExp :: Integer -> Integer-> Integer
1
2 fastExp _0 = 1
3 fastExp a 1 = a
4
  fastExp a b
5
     |b < 0 = undefined
6
     leven b = res*res
7
     |otherwise = let next=(fastExp a (b-1)) in (a * next)
8
      where res=(fastExp a (div b 2))
  Suppose we execute fastExp 2 7. The call stack would be
     fastExp 2 7 line 7 pattern match
     fastExp 2 6 line 6 pattern match
     fastExp 2 3 line 7 pattern match
     fastExp 2 2 line 6 pattern match
     fastExp 2 1 line 3 pattern match, STOP RECURSION
```

⁵Here for more informations: https://wiki.haskell.org/Let_vs_Where

```
Guards,let and where constructs
  fastExp :: Integer -> Integer-> Integer
1
2 fastExp _0 = 1
3 fastExp a 1 = a
4
  fastExp a b
5
     |b < 0 = undefined
6
     leven b = res*res
7
     |otherwise = let next=(fastExp a (b-1)) in (a * next)
8
      where res=(fastExp a (div b 2))
  Suppose we execute fastExp 2 7. The call stack would be
     fastExp 2 7 line 7 pattern match
     fastExp 2 6 line 6 pattern match
     fastExp 2 3 line 7 pattern match
     fastExp 2 2 line 6 pattern match
     fastExp 2 1 line 3 pattern match, STOP RECURSION
   In contrast to where, let are expressions and can be used
  anywhere<sup>5</sup>.
```

⁵Here for more informations: https://wiki.haskell.org/Let_vs_Where

If, case

```
if construct works as expected
```

```
1 div' n d = if d==0 then Nothing else Just (n/d)
```

```
case construct
```

Useful when we don't wish to define a function every time we need to do pattern matching.

```
f p11 ... p1k = e1
...
f pn1 ... pnk = en
--where each pij is a pattern,
--is semantically equivalent to:
f x1 x2 ... xk = case (x1, ..., xk) of
(p11, ..., p1k) -> e1
...
(pn1, ..., pnk) -> en
All patterns of a function return the same type hence all the
RHS of the case have the same type
```

case construct example

```
Pattern match "outside" the function definition. Note that all the
cases return the same type (a list of b's in this case)
cE :: (Ord a) :: a -> a -> [b]
cE a b xs = case (a 'compare' b,xs) of
        (_,[]) -> []
        (LT,xs) -> init xs
        (GT,xs) -> tail xs
        (E0,xs) -> [head xs]
```

ranges

Shortcut for listing stuff that can be enumerated. What if we need to test if a string contains a letter up to the lower case*j*? (Explicitly list all the letters is not the correct answer). ['a'..'j'] -- results in "abcdefghij" (String are [Char]) It work even in construction infinite list [1,3..] -- results in [1,3,5,7,9,11,13,15.....] and because of laziness we can (safely) do take 10 [1,3..]

- Collection of elements of the SAME TYPE.
- Delimited by square brackets and elements separated by commas.
- List che be consed. The cons operator (:) is used to incrementally build list putting an element at its head.
- empty list is []
- cons is a function that takes two parameter
 (:) :: a -> [a] -> [a]
 1:2:3:4:[]

list comprehension

It is a familiar concept for those who already have some experience in python It resambles the mathematical set specification. For instance let's compute the list of the factorial of the natural numbers

```
[product [2..x] | x<-[1..]]
```

```
More examples:
  [[2..x*2] | x<-[1..]]
  [filter (even) [2..x] | x<-[1..]]
  --:m Data.Char (ord)
  [let p=y*x in if even p then (negate p) else
   (p*2) |x<-[1..10], y<-(map ord ['a'..'z'])]
  --:m Data.List (nub)
nub $ map (\(x,y,z) -> z) [(a,b,c) | a<-[1..20],b<-[1..20],
   c<-[1..20], a^2+b^2==c^2, a+b+c>10]
```

Lambda functions - The Idea

Anonymous functions i.e. no need to give it a name

►
$$\lambda yx \rightarrow 2x + x^y$$
 translates in
(\x y -> 2*x + x^y)

- > Usually used withing high order function context. map (\x -> x*x-3) [1,10..300] map (\x -> let p = ord x in if even p then p else p^2) "Lambda functions are cool!"
- F = (\x₁..x_n−> exp(x₁..x_n))(v₁,..,v_k) substitute each occurence of the free variable x_i with the value v_i. If k < n f is again a function.</p>

Section 3

Coding - Problems on Lists

Given a polymorphic list *l* of type [*a*], find the last element of l (not using function *last*, l'm sorry).

Examples:

[1,2,3,4] = 4

Given a polymorphic list *l* of type [*a*], find the last element of l (not using function *last*, l'm sorry).

```
_last [1,2,3,4] = 4
_last ["programming","haskell","is","cool"]= "cool"
```

Given a polymorphic list *l* of type [*a*], find the last element of l (not using function *last*, l'm sorry).

```
_last [1,2,3,4] = 4
_last ["programming","haskell","is","cool"]= "cool"
```

Given a polymorphic list *I* of type [*a*], find the last element of I (not using function *last*, I'm sorry).

```
_last [1,2,3,4] = 4
_last ["programming","haskell","is","cool"]= "cool"
Solution
_last :: [a] -> a
_last [] = error "Undefined operation"
_last (x:[]) = x
_last (x:xs) = _last xs
```

Find the k'th element of a list where the first element has index 1

Examples:

elementAt 2 [3,35,32,33] = 35

Find the k'th element of a list where the first element has index 1

Examples:

elementAt 2 [3,35,32,33] = 35 elementAt 3 [('a',97),('b',98),('c',99)] = ('c',99) elementAt 4 [('a',97),('b',98),('c',99)] = error "Index out

Find the k'th element of a list where the first element has index 1

Examples:

elementAt 2 [3,35,32,33] = 35 elementAt 3 [('a',97),('b',98),('c',99)] = ('c',99) elementAt 4 [('a',97),('b',98),('c',99)] = error "Index out

Find the k'th element of a list where the first element has index 1

Examples:

```
elementAt 2 [3,35,32,33] = 35
elementAt 3 [('a',97),('b',98),('c',99)] = ('c',99)
elementAt 4 [('a',97),('b',98),('c',99)] = error "Index out
```

Solution

```
elementAt :: Integer -> [a] -> a
elementAt _ [] = error "index out of bound"
elementAt 1 (x:_) = x
elementAt n (_:xs) = elementAt (n-1) xs
```

Write a function that returns a boolean value tha indicates whether the input list is palindromic or not. $\ensuremath{\mathbf{1}}$

Write a function that returns a boolean value tha indicates whether the input list is palindromic or not. $1 \$

Examples:

palindrome "itopinonavevanonipoti" = True
palindrome "[1,2,3,3,1] = False

Write a function that returns a boolean value tha indicates whether the input list is palindromic or not. 1

Section 4

Problem on Numbers

Determine whether a given integer number is prime.

Determine whether a given integer number is prime.

Examples:

isPrime 57601 = True isPrime 1235 = False

Determine whether a given integer number is prime.

Examples: isPrime 57601 = True isPrime 1235 = False Solution isPrime n = _isPrime n 2 where __isPrime l k | k > l = True -- k > sqrt(l) | mod l k ==0 = False | otherwise = _isPrime l (k+1)

Implement the Euclid Method to find the greatest common divisor of two integer.

Implement the Euclid Method to find the greatest common divisor of two integer.

gcd '	30 12	=	6
gcd '	5 25	=	5
Implement the Euclid Method to find the greatest common divisor of two integer.

Examples:

gcd' 30 12 = 6 gcd' 5 25 = 5 Solution gcd' 0 y = y gcd' x y = gcd' (mod y x) x

Calculate Euler's totient function phi(m).

Euler's so-called totient function $\phi(m)$ is defined as the number of positive integers r ($1 \le r < m$) that are **coprime** to m.

Calculate Euler's totient function phi(m).

Euler's so-called totient function $\phi(m)$ is defined as the number of positive integers r ($1 \le r < m$) that are **coprime** to m.

Examples:

totient 10 = 4 totient 57601 = 57600 --57601 is prime^^

Calculate Euler's totient function phi(m).

Euler's so-called totient function $\phi(m)$ is defined as the number of positive integers r ($1 \le r < m$) that are **coprime** to m.

Examples:

totient 10 = 4 totient 57601 = 57600 --57601 is prime^^

Solution

```
totient n = length [e | e <- [1..n], coprime e n] where coprime e n = gcd n e ==1
```

Section 5

Find Best Variance - Stock Data

Write a program that read a file containing daily stock data. Each line of the file records data regarding prices of a good registered at regular time interval during each day. Fine the day which have the maximum variance between opening and closing price (second and last price record).

File content:

2012-03-30, **32.40**, 32.41, 32.04, 32.26, 31749400, **32.26** 2012-03-29, **32.06**, 32.19, 31.81, 32.12, 37038500, **32.12** 2012-03-28, **32.52**, 32.70, 32.04, 32.19, 41344800, **32.19**

Solution

```
The Solution. cabal install split
module Main where
import System.Environment (getArgs)
import Data.List.Split (splitOn)
import Data.List (maximumBy)
--main entry point
main = do
(fileName:_) <- getArgs</pre>
strF <- readFile fileName putStrLn $ maxDay strF</pre>
maxDay ::String -> String
maxDay s = snd $ maximum ss
 where
  ss = map (var . (splitOn ",")) $ lines s
var xs = abs diff
  where diff=((read (xs!!1)) - (read (last xs)), head xs)
```

Section 6

Coding - Project Euler Problem 1

If we list all the natural numbers below 10 that are multiples of 3 or 5, we get 3, 5, 6 and 9. The sum of these multiples is 23. Find the sum of all the multiples of 3 or 5 below 1000.

If we list all the natural numbers below 10 that are multiples of 3 or 5, we get 3, 5, 6 and 9. The sum of these multiples is 23. Find the sum of all the multiples of 3 or 5 below 1000.

How would you solve it using Haskell?

If we list all the natural numbers below 10 that are multiples of 3 or 5, we get 3, 5, 6 and 9. The sum of these multiples is 23. Find the sum of all the multiples of 3 or 5 below 1000.

How would you solve it using Haskell?

problem1' = sum .
 filter (\x -> x 'mod' 3==0 || x 'mod' 5 ==0)

Section 7

Coding - Project Euler Problem 26

A unit fraction contains 1 in the numerator. Where 0.1(6) means 0.1666666..., and has a 1-digit recurring cycle. It can be seen that 1/7 has a 6-digit recurring cycle. Find the value of d < 1000 for which 1/d contains the longest recurring cycle in its decimal fraction part.

- ▶ 1/2 = 0.5 0-recur
- ▶ 1/3 = 0.(3) 1-recur
- ▶ 1/4 = 0.25 0-recur
- ▶ 1/5 = 0.2 0-recur
- ▶ 1/6 = 0.1(6) 1-recur
- ▶ 1/7 = 0.(142857) 6-recur
- ▶ 1/8 = 0.125 0-recur
- ▶ 1/9 = 0.(1) 1-recur
- ▶ 1/10 = 0.1 0-recur

Key idea: Find the order of 10 in $\mathbb{N}/p\mathbb{N}$

The length of the repetend (period of the repeating decimal) of 1/p is equal to the order of 10 modulo p. If 10 is a primitive root modulo p, the repetend length is equal to p - 1; if not, the repetend length is a factor of p - 1. This result can be deduced from Fermat's little theorem, which states that $10p - 1 \equiv 1 \pmod{p}$. (Wikipedia)

The smallest power *n* of *g* s.t. $g^n \equiv 1 \pmod{p}$.

Problems 26 - Order finding example

- $10^{1} \equiv 10 \pmod{13}$ $10^{2} \equiv 9 \pmod{13}$ $10^{3} \equiv 12 \pmod{13}$ $10^{4} \equiv 3 \pmod{13}$ $10^{5} \equiv 4 \pmod{13}$ $10^{6} \equiv 1 \pmod{13}$
- ▶ 6 is the order of 10 (modulo 13)
 ▶ map (\a -> mod (10^a) 13) [1..12]

```
So now the problem is. Compute the order of numbers
n < 1000 and return the one that have maximum order
--modulo, current order
order :: Integer -> Integer -> Integer
order a ord
\mid mod (10^{ord}) a == 1 = ord
| \text{ord} > a = 0
otherwise = order a (ord+1)
maxo = fst $ maximumBy comparing $ pp
  where
    comparing = ((m,n) (p,q) \rightarrow n \text{ 'compare' } q)
     pp = map (\langle x - \rangle (x, order x 1))
           (filter (x \rightarrow mod x 10 > 0) [1,3..1000])
```

